AUTOMATING MOBILITY IN SMART CITIES

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2017 IFAC World Congress

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OUTLINE

What is a "Smart City" ?

Mobility as a resource contention game:

Selfish v Social optimality \Rightarrow PRICE OF ANARCHY (PoA)

Two takeaways (proposed research directions) from this talk:

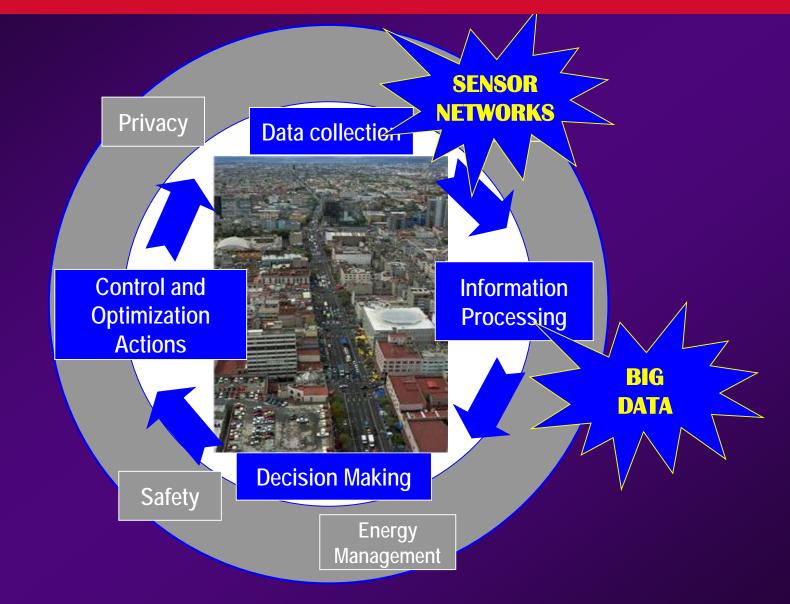
1. Use "Big Data" to estimate the PoA

2. Use Connected Autonomous Vehicles (CAVs) + control to reduce/eliminate the PoA

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"SMART CITY" AS A CYBER-PHYSICAL SYSTEM



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"SMART CITY" AS A CYBER-PHYSICAL SYSTEM



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WHAT IS A "SMART CITY" ?

"A city well performing in a forward-looking way in [economy, people, governance, mobility, environment, and living] built on the smart combination of endowments and activities of self-decisive, independent and aware citizens." *Giffinger et al, 2007*

Hitachi's vision for the Smart Sustainable City seeks to achieve concern for the global environment and lifestyle safety and convenience through the coordination of infrastructure. Smart Sustainable Cities realized through the coordination of infrastructures consist of two infrastructure layers that support consumers' lifestyles together with the urban management infrastructure that links these together using *Hitachi Web, 2014* Smart Sustainable Cities use information and communication technologies (ICT) to be more intelligent and efficient in the use of resources, resulting in cost and energy savings, improved service delivery and quality of life, and reduced environmental footprint--all supporting innovation and the low-carbon economy. *Cohen, 2014*

"We believe a city to be smart when investments in human and social capital and traditional (transport) and modern (ICT) communication infrastructure fuel sustainable economic growth and a high quality of life, with a wise management of natural resources, through participatory governance." *Meijer and Bolívar, 2013* Smart Sustainable Cities use information and communication technologies (ICT) to be more intelligent and efficient in the use of resources, resulting in cost and energy savings, improved service delivery and quality of life, and reduced environmental footprint--all supporting innovation and the low-carbon economy. *Cohen*, 2014

WHAT IS A "SMART CITY" ?

CREDIT: Fernando Livschitz

http://www.fastcodesign.com/3035870/filmmaker-creates-worlds-most-terrifying-traffic-intersection

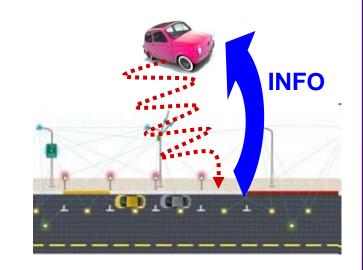
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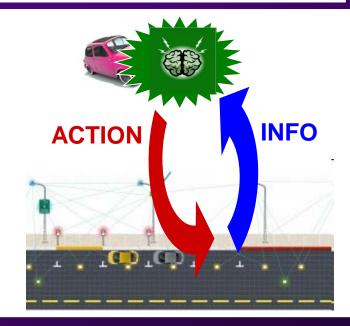
WHAT IS REALLY "SMART" ?

COLLECTING DATA IS NOT "SMART"

- JUST A NECESSARY STEP TO BEING "SMART"

PROCESSING DATA TO MAKE GOOD DECISIONS IS "SMART"





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TRAFFIC CONTROL



100-km Chinese traffic jam enters Day 9



The BU Bridge mess, Boston, MA (simulation using VISSIM)

WHY CAN'T WE IMPROVE TRAFFIC...

... EVEN IF WE KNOW THE ACHIEVABLE OPTIMUM IN A TRAFFIC NETWORK ???

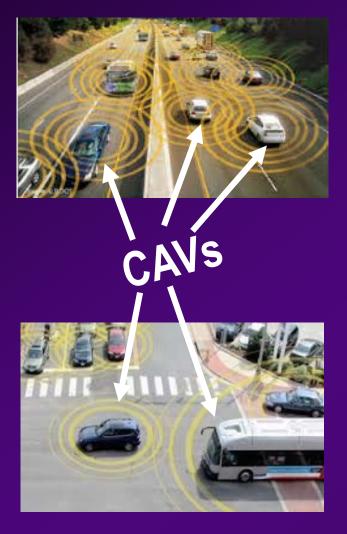
Because:

- Not enough controls (traffic lights, tolls, speed fines)
 → No chance to unleash the power of feedback!
- Not knowing other drivers' behavior leads to poor decisions (a simple game-theoretic fact)
 - → Drivers seek individual (selfish) optimum, not system-wide (social) optimum



PRICE OF ANARCHY (POA)

GAME-CHANGING OPPORTUNITY: CONNECTED AUTONOMOUS VEHICLES (CAVs)



FROM (SELFISH) "DRIVER OPTIMAL" TO (SOCIAL) "SYSTEM OPTIMAL" TRAFFIC CONTROL



NO TRAFFIC LIGHTS, NEVER STOP...

THE CASE FOR "SELF-DRIVING" CARS

- Humans are bad drivers (94% of accidents are due to human error)
- Computers do not get distracted (humans do)
- Computers can process vast amounts of data (humans cannot)
- Computers can maintain steady cruising speeds (leading to improved energy efficiency)
- Computers react quickly (humans do not)
- Computer can make fast and accurate driving adjustments

Computers do not blink, do not drink, and do not sleep

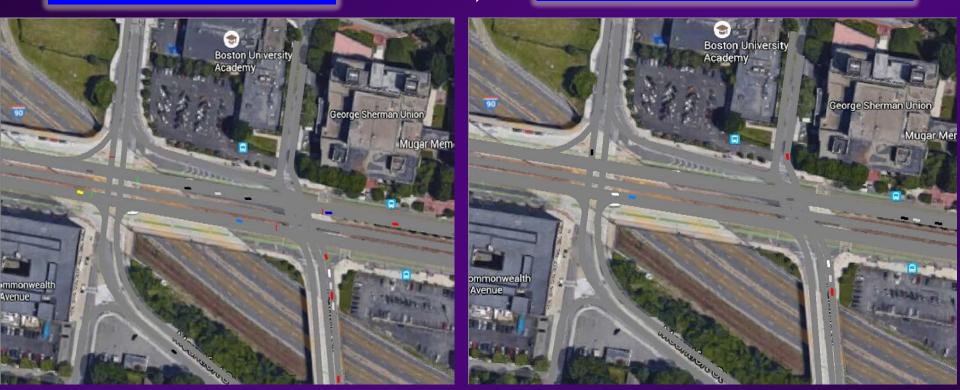
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- Legal issues: who is to blame for a fault ?
- How to integrate with normal cars ?
- Security and Privacy (due to connectivity)
- Accidents may be rare, but when they occur they are likely to be serious
- Technical challenges...

HOW TO QUANTIFY BENEFITS OF AUTOMATED MOBILITY ?

BEFORE (traffic lights)

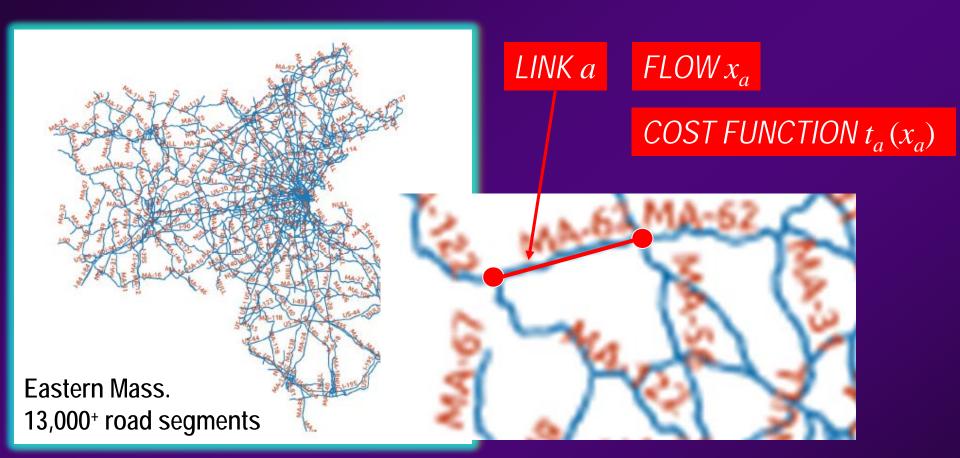
AFTER (no traffic lights, CAVs)



One of the worst-designed double intersections ever... (BU Bridge – Commonwealth Ave, Boston, MA)

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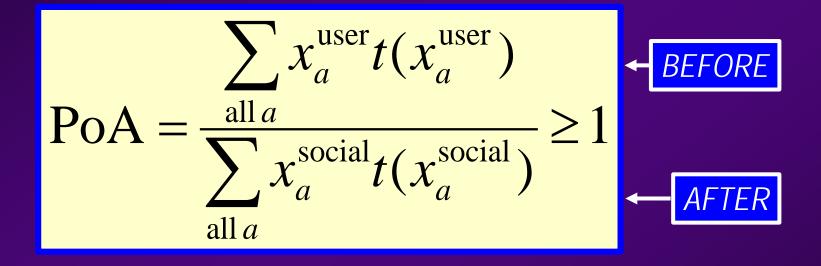
HOW TO MEASURE THE PRICE OF ANARCHY ?



Under USER-CENTRIC (selfish) control: x_a^{user} is the equilibrium flow Under SYSTEM-CENTRIC (social) control: x_a^{social} is the equilibrium flow

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HOW TO MEASURE THE PRICE OF ANARCHY ?



Two takeaways (proposed research directions) from this talk:

- 1. Measure/estimate the PoA?
- 2. Reduce/eliminate the PoA?

Inverse Optimization

Optimal Control Framework

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ESTIMATING THE PRICE OF ANARCHY



CHALLENGES AS THINGS NOW STAND...

- We don't know user COST FUNCTIONS
- We don't know user ORIGIN-DESTINATION pairs (no DEMAND model)



We can't exploit CAVs

We can't assess the value of investing in CAV-based technologies, since we can't evaluate the PRICE OF ANARCHY

BUT WE DO HAVE PLENTY OF DATA....

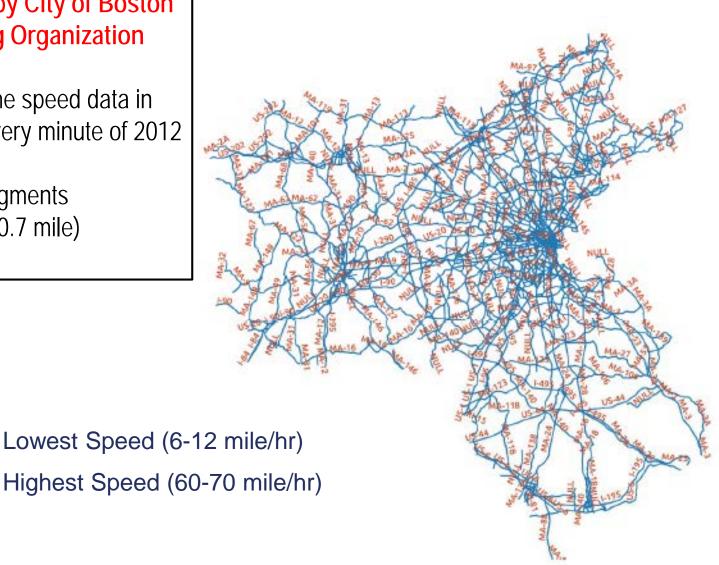
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THE BIG DATA SET

Data Set provided by City of Boston and Mass. Planning Organization

Average and real-time speed data in Eastern Mass. for every minute of 2012

- 13000⁺ road segments (avg. distance 0.7 mile)
- 50⁺ GB of data



INVERSE OPTIMIZATION – KEY IDEA

BACKWARD optimization:

- Data reveal a (selfish) equilibrium (Wardrop/Nash equilibrium)
- What are the (virtual) cost functions which best fit the data that lead to this equilibrium?



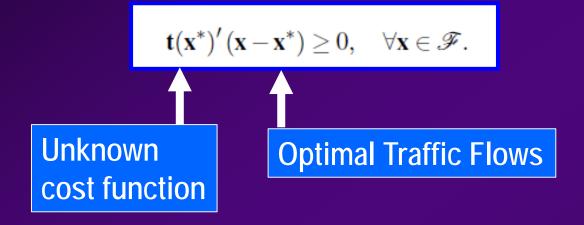
FORWARD optimization:

- Use these cost functions to find (social) optimal traffic flows

- Estimate the PRICE OF ANARCHY

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Optimal Traffic Flow allocation as a Variational Inequality (VI) problem:



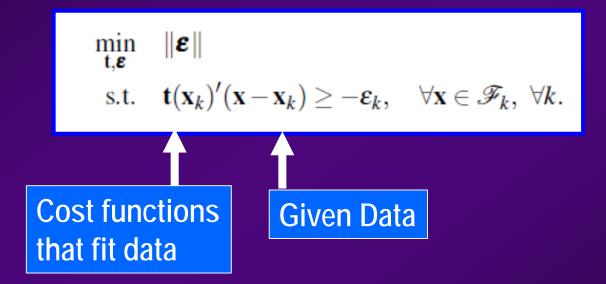
Assumption 1: $t(\cdot)$ is strongly monotone and continuously differentiable. \mathscr{F} is nonempty and contains an interior point (Slater's condition)

THEOREM: Suppose Assumption 1 holds. Then, there exists a Wardrop (Nash) equilibrium of the single-class transportation network which is the unique solution of the VI problem.

Patriksson, 2015

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Inverse Variational inequality problem



Solve for cost functions $t(\cdot)$ in a *Reproducing Kernel Hilbert Space*

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Data reconciliation

$$\begin{array}{c} \min_{f,\mathbf{y},\boldsymbol{\epsilon}} & \|\boldsymbol{\epsilon}\| + \gamma \|f\|_{\mathcal{H}}^{2} \\
\text{s.t.} \quad \mathbf{e}_{a}' \mathbf{N}_{k}' \mathbf{y}^{\mathbf{w}} \leq t_{a}^{0} f\left(\frac{x_{a}}{m_{a}}\right), \\
\forall \mathbf{w} \in \mathcal{W}^{(k)}, \ a \in \mathcal{A}^{(k)}, \ k \in \llbracket \mathcal{K} \rrbracket, \\
\sum_{a \in \mathcal{A}^{(k)}} t_{a}^{0} x_{a} f\left(\frac{x_{a}}{m_{a}}\right) - \sum_{\mathbf{w} \in \mathcal{W}^{(k)}} (\mathbf{d}^{\mathbf{w}})' \mathbf{y}^{\mathbf{w}} \leq \epsilon_{k}, \\
\forall k \in \llbracket \mathcal{K} \rrbracket, \\
f\left(\frac{x_{a}}{m_{a}}\right) \leq f\left(\frac{x_{\tilde{a}}}{m_{\tilde{a}}}\right), \\
\forall a, \ \tilde{a} \in \bigcup_{k=1}^{|\mathcal{K}|} \mathcal{A}^{(k)} \text{ s.t. } \frac{x_{a}}{m_{a}} \leq \frac{x_{\tilde{a}}}{m_{\tilde{a}}}, \\
\epsilon \geq \mathbf{0}, \quad f \in \mathcal{H}, \\
f(0) = 1,
\end{array}$$

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n

Reproducing Kernel:
$$\phi(x, y) = (c + xy)^n = \sum_{i=0}^n \binom{n}{i} c^{n-i} x^i y^i$$

$$\lim_{\substack{\beta, y, \epsilon}} \|\epsilon\| + \gamma \sum_{i=0}^n \frac{\beta_i^2}{\binom{n}{i} c^{n-i}}$$
s.t. $\mathbf{e}'_a \mathbf{N}'_k \mathbf{y}^{\mathbf{w}} \le t_a^0 \sum_{i=0}^n \beta_i \left(\frac{x_a}{m_a}\right)^i$,
$$\forall \mathbf{w} \in \mathcal{W}^{(k)}, \ a \in \mathcal{A}^{(k)}, \ k \in \llbracket \mathcal{K} \rrbracket,$$

$$\sum_{a \in \mathcal{A}_k} t_a^0 x_a \sum_{i=0}^n \beta_i \left(\frac{x_a}{m_a}\right)^i - \sum_{\mathbf{w} \in \mathcal{W}_k} (\mathbf{d}^{\mathbf{w}})' \mathbf{y}^{\mathbf{w}} \le \epsilon_k,$$

$$\forall k \in \llbracket \mathcal{K} \rrbracket,$$

$$\sum_{i=0}^n \beta_i \left(\frac{x_a}{m_a}\right)^i \le \sum_{i=0}^n \beta_i \left(\frac{x_{\tilde{a}}}{m_{\tilde{a}}}\right)^i,$$

$$\forall a, \ \tilde{a} \in \bigcup_{k=1}^{|\mathcal{K}|} \mathcal{A}^{(k)} \text{ s.t. } \frac{x_a}{m_a} \le \frac{x_{\tilde{a}}}{m_{\tilde{a}}},$$

$$\epsilon \ge 0, \quad \beta_0 = 1,$$

Link Cost Function Estimates

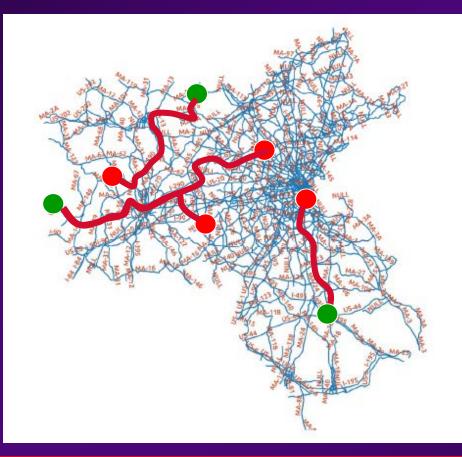
$$\hat{f}(x) = 1 + \sum_{i=1}^{n} \beta_i^* x^i$$

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O-D MATRIX ESTIMATION

That's a separate challenging problem!

We solve it using Generalized Least Squares (GLS) methods

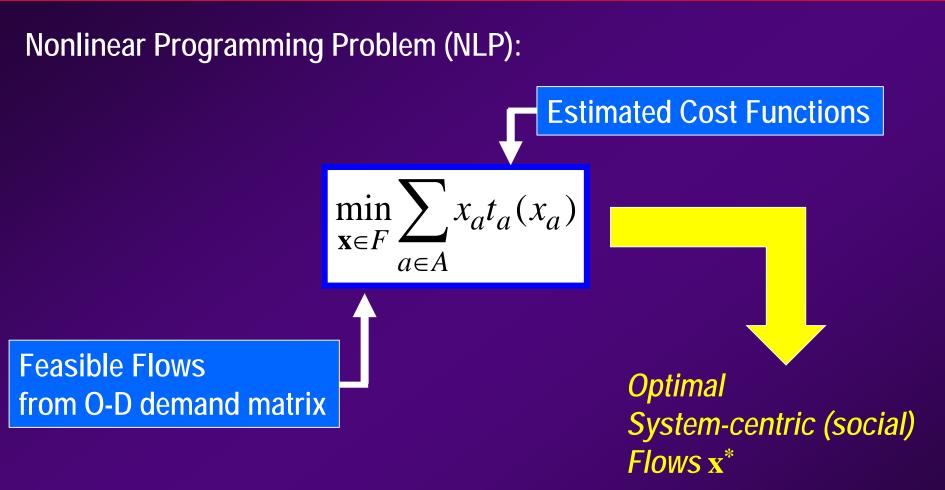


Zhang et al, IFAC 2017 - WeP13.1

Typically, hundreds/thousands of O-D pairs...

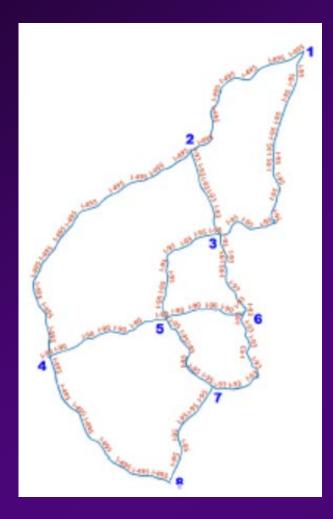
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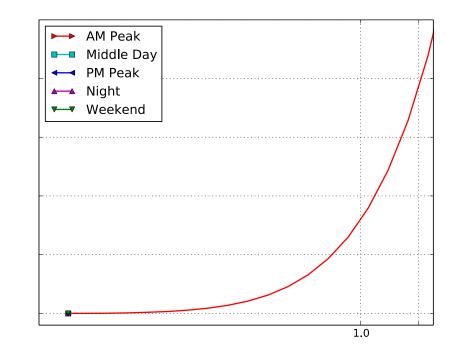
FORWARD OPTIMIZATION PROBLEM



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COST FUNCTION ESTIMATES: BOSTON AREA 2012

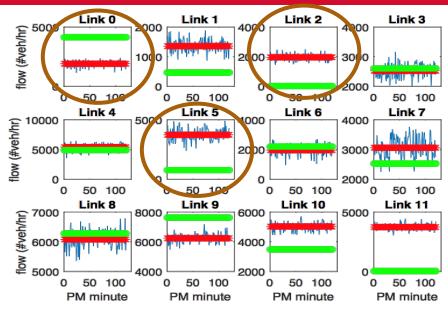


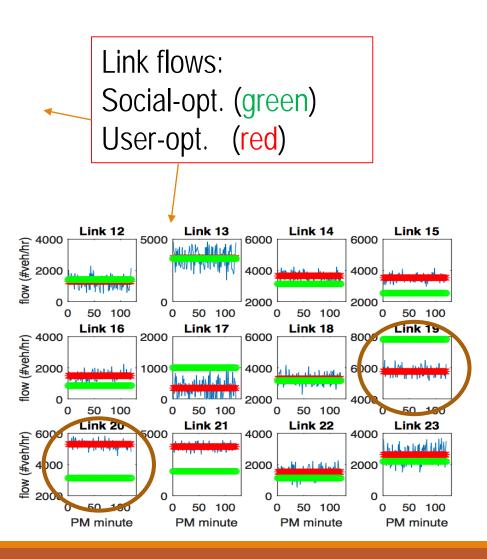


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PRICE OF ANARCHY – BOSTON AREA 2012

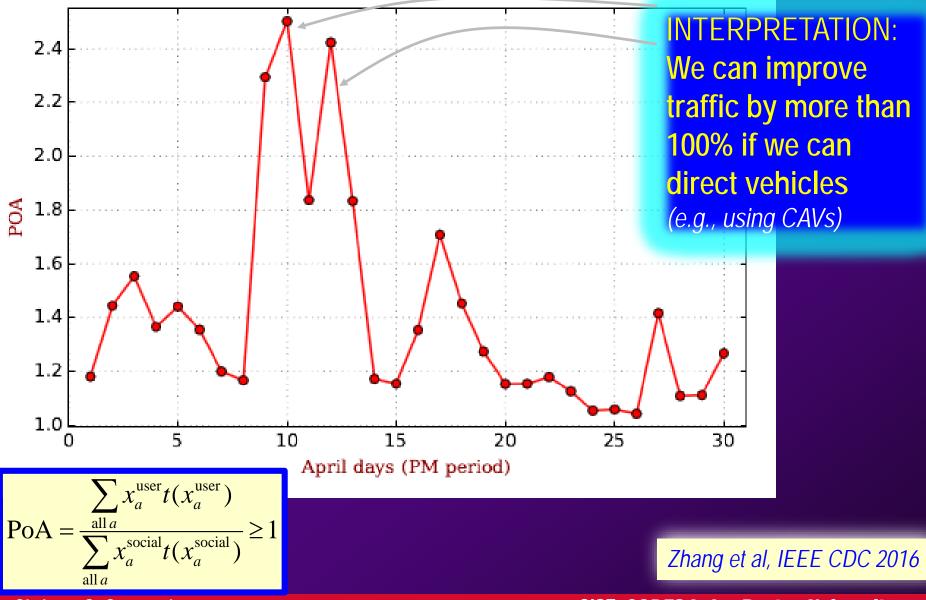




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PRICE OF ANARCHY – BOSTON AREA 2012



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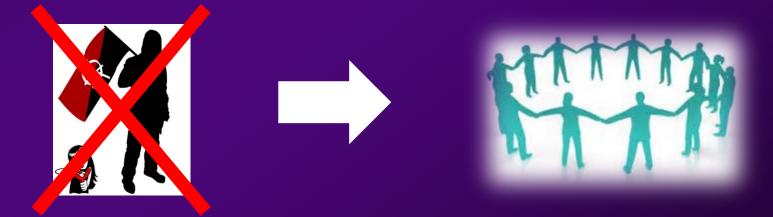
NEXT STEPS...

Goal 1 accomplished:

PoA is HIGH

⇒ Evidence to support investing in CAVs to achieve System-Centric (Social) Optimality

 \Rightarrow How do we do it ?



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Relatively easy:

- Recommendations/Suggestions (e.g., through apps)
- Incentives, Pricing Schemes
- Automated Enforcement (CAV selects route, driver can override)

NEXT STEPS – TRAFFIC NETWORK BOTTLENECKS

Need to automatically control velocity/acceleration in urban environments:

- Merging points
- Intersections



A DECENTRALIZED OPTIMAL CONTROL FRAMEWORK FOR CAVS



NO TRAFFIC LIGHTS, NEVER STOP...

VEHICLE COORDINATION

- Centralized approaches:
- Reservation schemes:

Dresner and Stone (2004), Huang et al (2012), Zhang et al (2013), Kim and Kumar (2014), Zhu and Ukkusuri(2015)

Control and Optimization:

Levine and Athans (1966), Varaiya (1993), Lu and Hedrick (2000), Kotsialos and Papageorgiou (2004), Li and Wang (2006), Lee and Park (2013), Kamal et al (2013), Pasquale et al (2015)

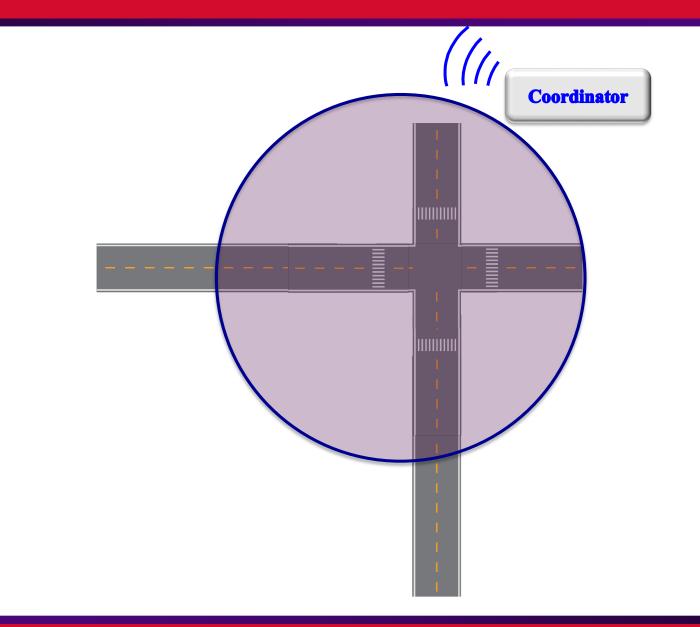
- Queueing models: Miculescu and Karaman (2014)
- Decentralized approaches:
 - Heuristics: Milanes et al (2011), Onieva et al (2012)
 - Critical set: Hafner et al (2013), Colombo and Del Vecchio (2014)
 - Optimization: Makarem et al (2013), Campos et al (2014)

RELATED RESEARCH

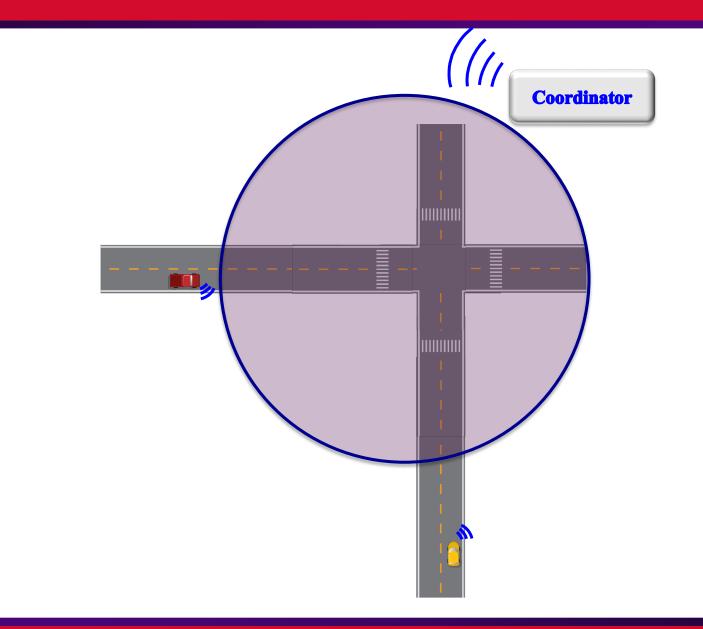
- KTH Royal Inst. of Tech., Sweden (K. Johansson et al)
- GIPSA-Lab, Grenoble, France (C. Canudas de Wit)
- U. of Genova, U. of Pavia (S. Sacone, S. Siri, A. Ferrara et al)
- Nanyang Technological University, Singapore (R. Su et al)
- Tsinghua U., China (Y. Zhang)
- MIT (A. Annaswamy, S. Karaman et al)
- UC Berkeley (A. Bayen et al)
- U. Michigan/Mcity (H. Peng et al)

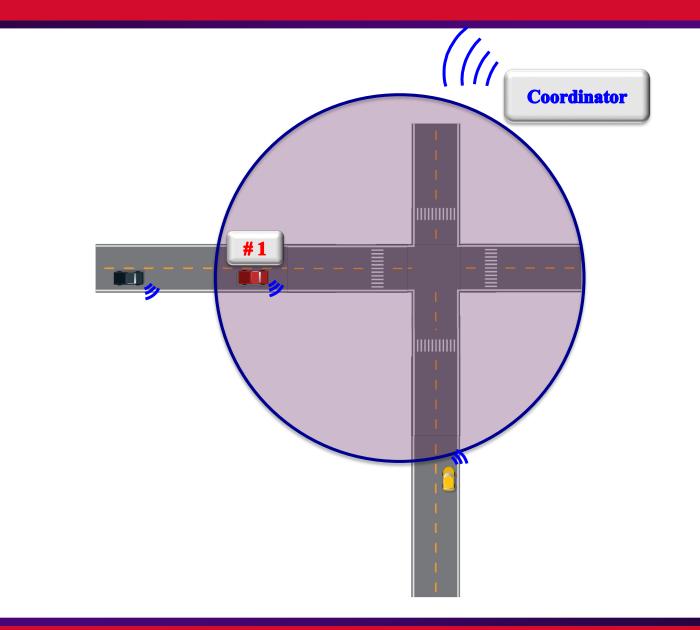
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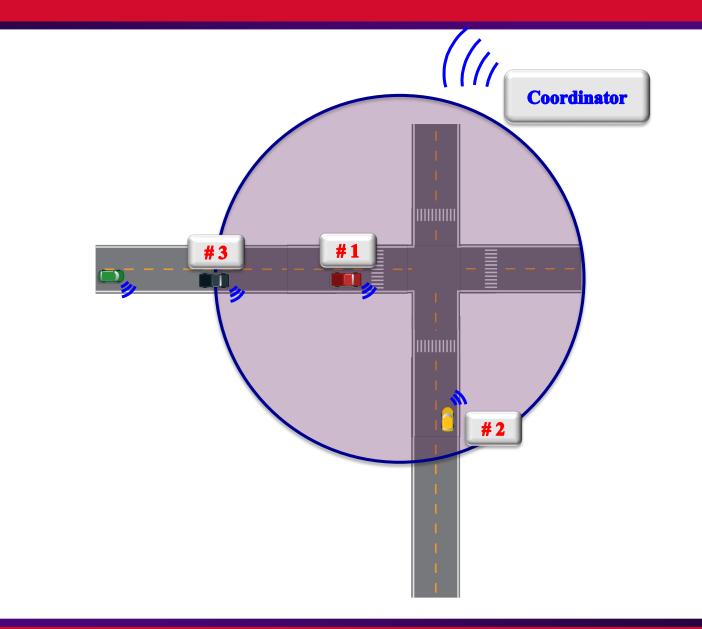
NO TRAFFIC LIGHTS - CAVs

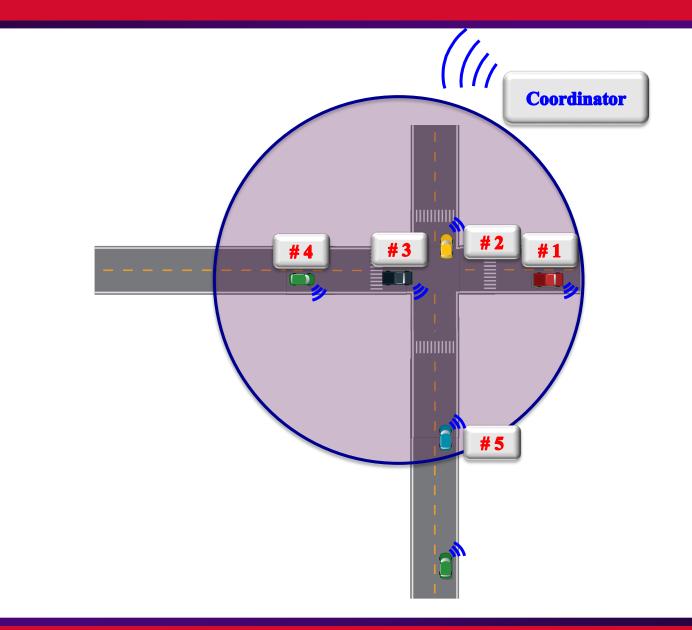


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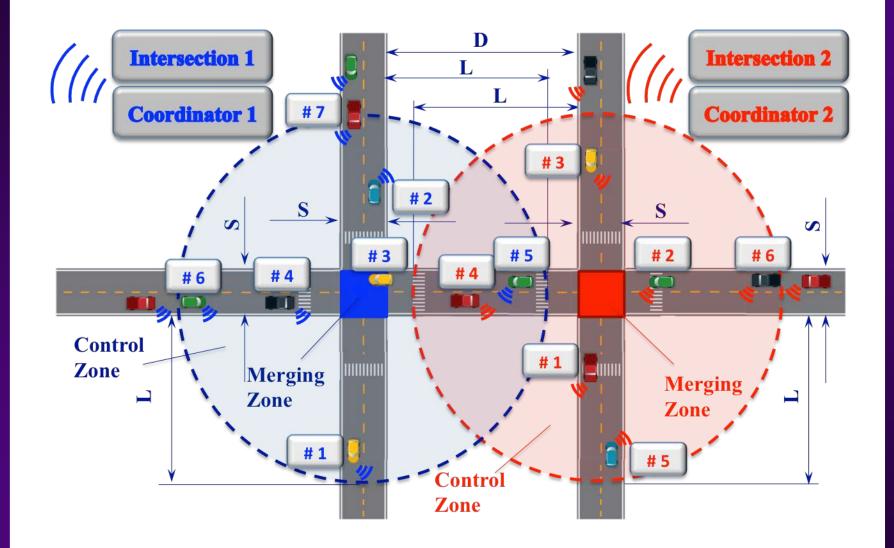






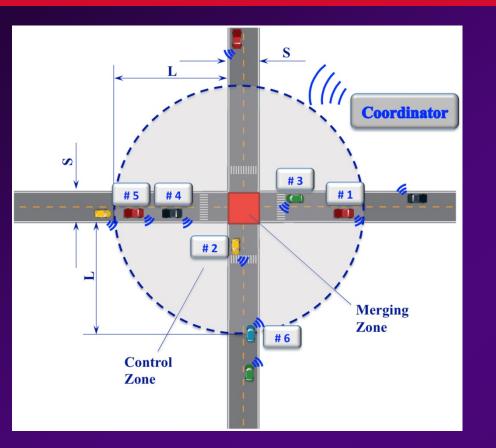


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THE MODEL



CAV dynamics:

$$\dot{p}_i = v_i(t)$$
$$\dot{v}_i = u_i(t)$$
$$t \in [t_i^0, t_i^f]$$

$$t_i^0$$
: Enters Control Zone (CZ)

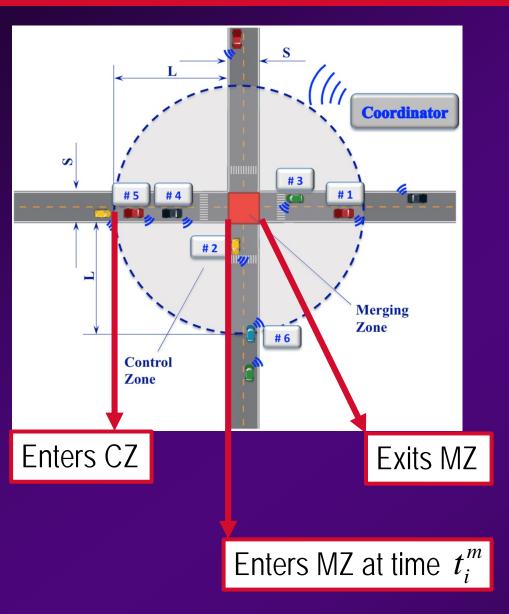
$$t_i^f$$
: Exits Merging Zone (MZ)

Speed, Acceleration constraints:

$$u_{\min} \le u_i(t) \le u_{\max}$$
$$0 \le v_{\min} \le v_i(t) \le v_{\max}$$

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THE MODEL



Control Zone queue:

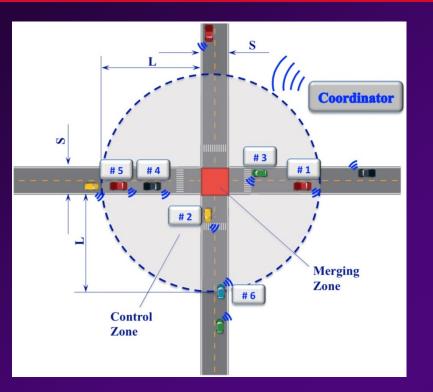
$$\mathcal{N}(t) = \{1, \dots, N(t)\}$$

Order constraint:

$$t_i^m \ge t_i^{m-1}, i \in \mathcal{N}(t), i > 1$$

Not necessarily FIFO – can change order at CAV arrival events

THE MODEL



Depending on physical location of *i* relative to *i*-1, *i*-1 belongs to one of the four subsets:

1. $R_i(t)$: same road, same direction as i, different lanes

2. $L_i(t)$: same road, same lane as i

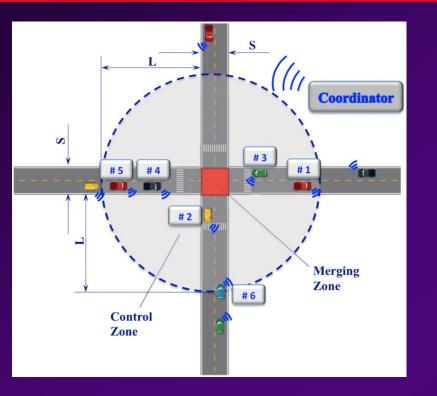
3. $C_i(t)$: different road from i,

possible collision at MZ

4. $O_i(t)$: same road as *i*, opposite direction, no collision at MZ

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THE MODEL - ASSUMPTIONS



No turn or lane change

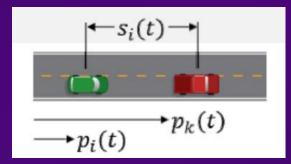
Constant speed in MZ:

$$t_i^f = t_i^m + \frac{S}{v_i(t_i^m)}$$

SAFETY CONSTRAINTS

Rear end safety constraint:

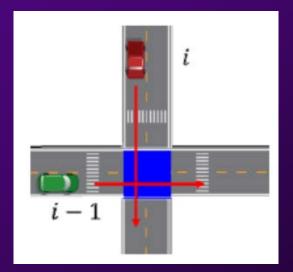
$$s_i(t) = p_k(t) - p_i(t) \ge \delta, \ t \in [t_i^0, t_i^f]$$



Lateral collision avoidance constraint:

$$\Gamma_{i} = \left\{ t : t \in [t_{i}^{m}, t_{i}^{f}] \right\}$$

$$\Gamma_{i} \cap \Gamma_{j} = \emptyset, \ t \in [t_{i}^{m}, t_{i}^{f}], \ j \in C_{i}(t)$$



ENERGY MINIMIZATION PROBLEM: *E-MIN*

$$\min_{u_i(t)} J_i(u_i(t), t_i^m) = \frac{1}{2} \int_{t_i^0}^{t_i^m} u_i^2(t) dt$$

subject to: 1. CAV dynamics

- 2. Speed/Acceleration constraints
- 3. Order constraints: $t_i^m \ge t_i^{m-1}$
- 4. Rear-end safety constraint
- 5. Lateral collision avoidance constraint

$$p_i(t_i^0) = 0, \quad p_i(t_i^m) = L$$

Given $t_i^0, v_i(t_i^0), \quad t_i^m \longrightarrow$ How is this determined?

Each CAV minimizes ENERGY COST FUNCTIONAL

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ENERGY MINIMIZATION PROBLEM: E-MIN

Feasible control set for *E-MIN*:

 $A_i = \{u_i(t) \in U_i \text{ subject to}:$

- 1. CAV dynamics
- 2. Speed/Acceleration constraints
- 3. Order constraints: $t_i^m \ge t_i^{m-1}$
- 4. Rear-end safety constraint
- 5. Lateral collision avoidance constraint $p_i(t_i^0) = 0, \quad p_i(t_i^m) = L$ Given $t_i^0, v_i(t_i^0), t_i^m$

HOW IS *i*th MERGING TIME DETERMINED ?

Maximize THROUGHPUT – Problem TP-MAX

$$\begin{split} \min_{\mathbf{t}_{(2:N(t))}} & \sum_{i=2}^{N(t)} \left(t_{i}^{m}(\mathbf{u}_{(1:i)}(t)) - t_{i-1}^{m}(\mathbf{u}_{(1:i-1)}(t)) \right) \\ &= \min_{\mathbf{t}_{N(t)}} \left(t_{N(t)}^{m}(\mathbf{u}_{(1:i)}(t)) - t_{1}^{m}(\mathbf{u}_{(1)}(t)) \right) \\ &\text{subject to : } u_{i}(t;t_{i}^{m}) \in \mathcal{A}_{i}, i \in \mathcal{N}(t) \\ &\quad s_{i}(t) = p_{k}(t) - p_{i}(t) \geq \delta, t \in [t_{i}^{0}, t_{i}^{m}] \\ &\quad t_{i}^{m} \geq t_{i-1}^{m}, i \in \mathcal{N}(t), i > 1 \end{split}$$

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HOW IS *i*th MERGING TIME DETERMINED ?

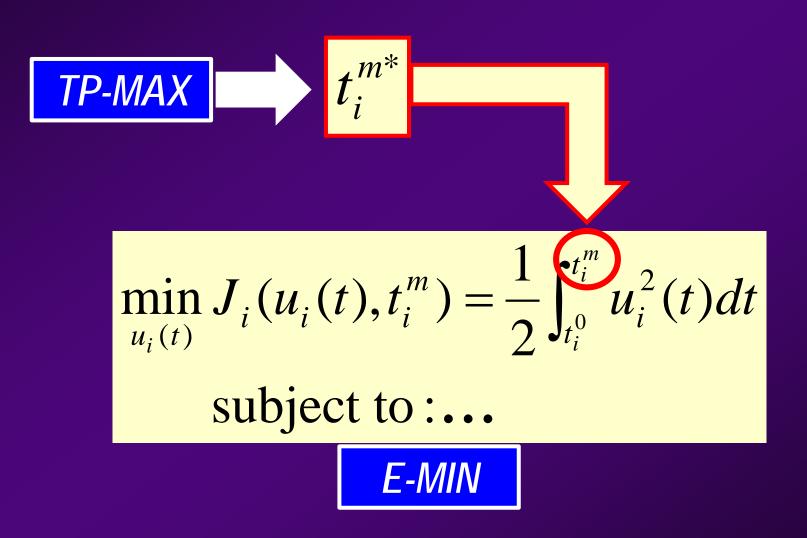
THEOREM:

The solution of TP-MAX is recursively determined by each *i*:

$$t_{i}^{m^{*}} = \begin{cases} t_{1}^{m^{*}} & \text{if } i = 1 \\ \max\{t_{i-1}^{m^{*}}, t_{k}^{m^{*}} + \frac{\delta}{v_{k}^{m}}, t_{i}^{c}\} & \text{if } i - 1 \in R_{i}(t) \cup O_{i}(t) \\ \max\{t_{i-1}^{m^{*}} + \frac{\delta}{v_{i-1}^{m}}, t_{i}^{c}\} & \text{if } i - 1 \in L_{i}(t) \\ \max\{t_{i-1}^{m^{*}} + \frac{S}{v_{i-1}^{m}}, t_{i}^{c}\} & \text{if } i - 1 \in C_{i}(t) \end{cases}$$
where: $t_{i}^{c} = t_{i}^{1} \mathbf{1}_{v_{i}^{m} = v_{\max}} + t_{i}^{2}(1 - \mathbf{1}_{v_{i}^{m} = v_{\max}}) \quad t_{i}^{1} = t_{i}^{0} + \frac{L}{v_{\max}} + \frac{(v_{\max} - v_{i}^{0})^{2}}{2u_{i,\max}v_{\max}} \\ t_{i}^{2} = t_{i}^{0} + \frac{[2Lu_{i,\max} + (v_{i}^{0})^{2})^{1/2} - v_{i}^{0}}{2u_{i,\max}} \end{cases}$
Known constant

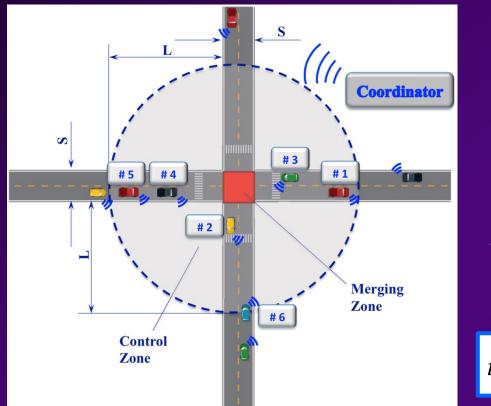
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HOW IS *i*th MERGING TIME DETERMINED ?



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DECENTRALIZED FRAMEWORK



CAV INFORMATION SET upon entering a CZ:

$$Y_{i}(t) = \left\{ p_{i}(t), v_{i}(t), w, Q_{i}, s_{i}(t), t_{i}^{m^{*}} \right\}$$

w : unique CAV ID Q_i : one of the four sets R_i , L_i , C_i , O_i

 $t_i^{m^*}$: solution of TP-MAX obtained by *i*

INFORMATION SET available to *i* and COORDINATOR upon entering CZ
 Communication needed (e.g., DSDC)

Communication needed (e.g., DSRC)

DECENTRALIZED PROBLEM FOR EACH CAV *i*

$$\min_{u_i(t)} \frac{1}{2} \int_{t_i^0}^{t_i^{m^*}} u_i^2(t) dt$$

subject to: 1. CAV dynamics
2. Speed/Acceleration constraints
3. $t_i^{m^*}$ from TP-MAX solution
 $p_i(t_i^0) = 0, \quad p_i(t_i^m) = L$
Given t_i^0, v_i^0

NOT INCLUDED:

Lateral collision avoidance constraint

Implicitly handled by $t_i^{m^*}$

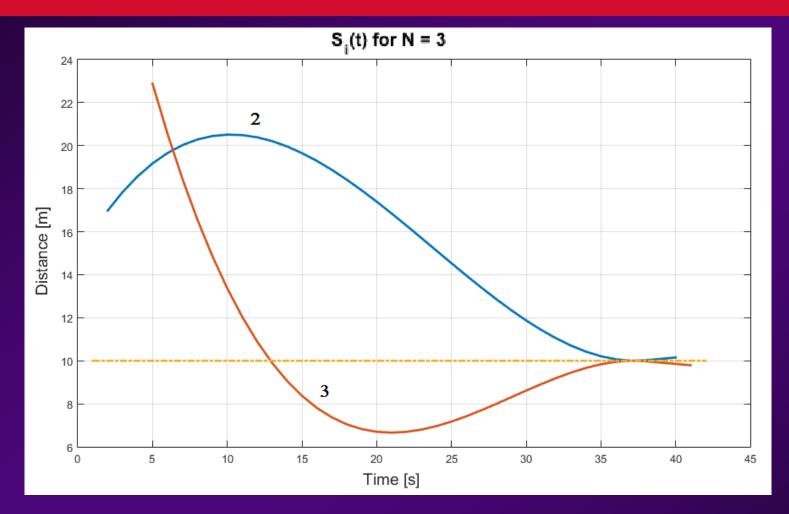
Rear-end safety constraint

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Only guaranteed at $t_i^{m^*}$

SAFETY CONSTRAINT NOT GUARANTEED...



Safety constraint violation by CAV 3 when $\delta = 10$.

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DECENTRALIZED PROBLEM SOLUTION

When constraints are not active:

$$u_i^*(t) = a_i t + b_i$$

$$v_i^*(t) = \frac{1}{2}a_it^2 + b_it + c_i$$

$$p_i^*(t) = \frac{1}{6}a_i t^3 + \frac{1}{2}b_i t^2 + c_i t + d_i$$

Coefficients obtained from:

$$\begin{bmatrix} \frac{1}{6}(t_i^0)^3 & \frac{1}{2}(t_i^0)^2 & t_i^0 & 1\\ \frac{1}{2}(t_i^0)^2 & t_i^0 & 1 & 0\\ \frac{1}{6}(t_i^m)^3 & \frac{1}{2}(t_i^m)^2 & t_i^m & 1\\ -t_i^m & -1 & 0 & 0 \end{bmatrix} \cdot \begin{bmatrix} a_i \\ b_i \\ b_i \\ c_i \\ d_i \end{bmatrix} = \begin{bmatrix} p_i(t_i^0) \\ v_i(t_i^0) \\ p_i(t_i^m) \\ 0 \end{bmatrix}$$

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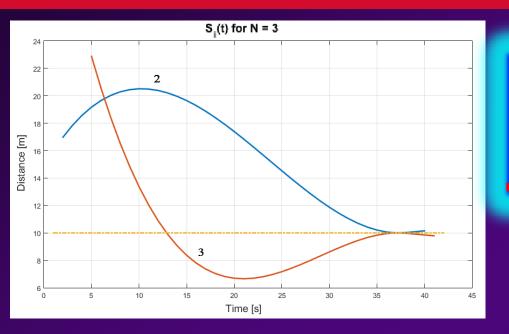
DECENTRALIZED PROBLEM SOLUTION

When one or more constraints are active:

Solution is of the same form and still analytically tractable

Malikopoulos, Cassandras, and Zhang, 2017

FEASIBILITY ANALYSIS



Under what conditions can we guarantee safety throughout the CZ ?

THEOREM:

There exists a nonempty feasible region of initial conditions (t_i^0, v_i^0) for each *i* such that, under the decentralized optimal control, $s_i(t) = p_k(t) - p_i(t) \ge \delta$ holds for all $t \in [t_i^0, t_i^m]$ given initial conditions $t_k^0, v_k^0, t_k^m, v_k^0$ for *k*

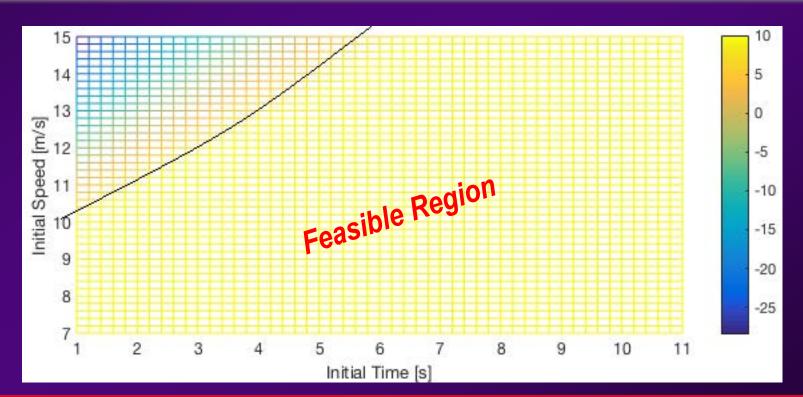
Zhang, Cassandras, and Malikopoulos, ACC 2017

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FEASIBILITY ANALYSIS

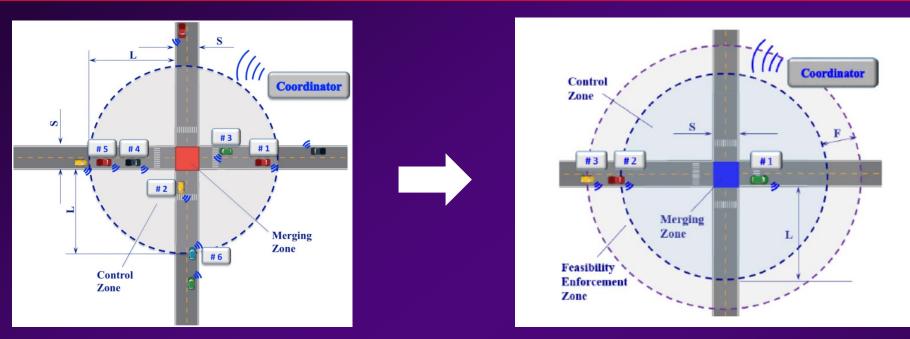
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FEASIBILITY ENFORCEMENT ZONE



What is the length of the Feasibility Enforcement Zone (FEZ) ? Worst case analysis:

When CAV *i* enters FEZ with v_{max} and needs to reach CZ with v_{min}

$$\overrightarrow{F} = \frac{v_{\min}^2 - v_{\max}^2}{2u_B}, \quad u_B : \text{min. acceleration s.t. } u_{\min} < u_B < 0$$

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WHO NEEDS TRAFFIC LIGHTS?

With traffic lights

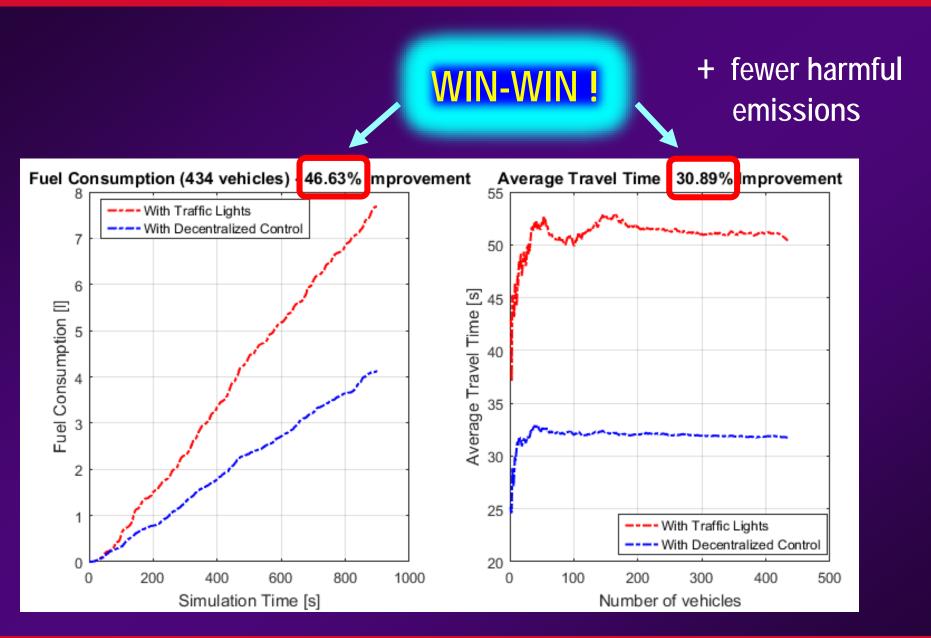
With decentralized control of CAVs



One of the worst-designed double intersections ever... (BU Bridge – Commonwealth Ave, Boston, MA)

Christos G. Cassandras

EXAMPLE



Christos G. Cassandras

TESTING AUTOMATED MOBILITY

BU Robotics Lab

Mcity test bed, U. Michigan



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OPTIMAL CONTROL FRAMEWORK – ISSUES

- Computational complexity for on-line implementation
- Incorporating turns and extending to multiple intersections
- Alternative formulations: travel time + fuel efficiency
- How about pedestrians and non-CAV traffic?

CONCLUSIONS

Two takeaways (proposed research directions) from this talk:

- 1. Use real data to infer user behavior and solve systemcentric problems, estimate Price of Anarchy (PoA)
- 2. Use Connected Autonomous Vehicles (CAVs) + control to reduce/eliminate the PoA

Interesting OPEN QUESTIONS regarding Automated Mobility:

- What fraction of CAVs does it take to realize benefits ?
- How do we integrate CAVs with pedestrians, bicycles, etc ?
- Is Shared Mobility On-Demand the long-term answer ? (typical car utilization is 4%...)

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Thank you Merci...